

# A Geometrical Problem Arising in a Signal Restoration Algorithm

Nir Cohen and Roy Meshulam

Communicated by Todd Quinto

**ABSTRACT.** Let  $\mathbb{Z}_{2N} = \{0, \dots, 2N - 1\}$  denote the group of integers modulo  $2N$ , and let  $L$  be the space of all real functions on  $\mathbb{Z}_{2N}$  which are supported on  $\{0, \dots, N - 1\}$ . The spectral phase of a function  $f : \mathbb{Z}_{2N} \rightarrow \mathbb{R}$  is given by  $\phi_f(k) = \arg \widehat{f}(k)$  for  $k \in \mathbb{Z}_{2N}$ , where  $\widehat{f}$  denotes the discrete Fourier transform of  $f$ .

For a fixed  $s \in L$  let  $K_s$  denote the cone of all  $f : \mathbb{Z}_{2N} \rightarrow \mathbb{R}$  which satisfy  $\phi_f \equiv \phi_s$ , and let  $M_s$  be its linear span. The angle  $\alpha_s$  between  $M_s$  and  $L$  determines the convergence rate of the signal restoration from phase algorithm of Levi and Stark [3]. Here we prove the following conjectures of Urieli et al. [7] who verified them for the  $N \leq 3$  case:

1.  $\alpha(M_s, L) \leq \pi/4$  for a generic  $s \in L$ .
2. If  $s \in L$  is geometric, i.e.,  $s(j) = q^j$  for  $0 \leq j \leq N - 1$  where  $\pm 1 \neq q \in \mathbb{R}$ , then  $\alpha(M_s, L) = \pi/4$ .

## 1. Introduction

Let  $\mathbb{Z}_{2N}$  denote the group of integers modulo  $2N$  ( $N \geq 2$ ), and let  $L_{\mathbb{C}}(\mathbb{Z}_{2N})$  ( $L_{\mathbb{R}}(\mathbb{Z}_{2N})$ ) denote the space of complex (real) valued functions on  $\mathbb{Z}_{2N}$  with the usual inner product  $\langle f, g \rangle = \sum_{x \in \mathbb{Z}_{2N}} f(x)\overline{g(x)}$ . The Fourier transform  $F : L_{\mathbb{C}}(\mathbb{Z}_{2N}) \rightarrow L_{\mathbb{C}}(\mathbb{Z}_{2N})$  is given by  $F(f)(y) = \widehat{f}(y) = \sum_{x \in \mathbb{Z}_{2N}} f(x)\omega^{-xy}$  where  $\omega = \exp(\pi i/N)$ . The spectral phase of  $f \in L_{\mathbb{R}}(\mathbb{Z}_{2N})$  is given by  $\phi_f(k) = \arg \widehat{f}(k)$  for  $k \in \mathbb{Z}_{2N}$ .

A natural and important problem which arises in diverse applications is the reconstruction of a Fourier transform pair  $f, F(f)$  from partial data on either or both functions. See Barakat and Newsam's papers [1, 2] for a study of various reconstruction algorithms and their analysis. One particular instance of the general reconstruction problem which is commonly encountered in image processing is the retrieval of a signal  $f \in L_{\mathbb{R}}(\mathbb{Z}_{2N})$  from its spectral phase  $\phi_f$ .

*Math Subject Classifications.* 94A11, 65T05.

*Keywords and Phrases.* Discrete Fourier Transform, Fourier pair reconstruction, Hermitian forms.

*Acknowledgments and Notes.* Nir Cohen – Supported by CNPq grant 300019/96-3.

Roy Meshulam – Research supported by the Fund for the Promotion of Research at the Technion.

Denote  $\text{Supp } f = \{x \in \mathbb{Z}_{2N} : f(x) \neq 0\}$ , and let  $L = \{s \in L_{\mathbb{R}}(\mathbb{Z}_{2N}) : \text{Supp } s \subset \{0, \dots, N - 1\}\}$ . For a function  $s \in L$ , let  $K_s$  denote the cone of all  $f \in L_{\mathbb{R}}(\mathbb{Z}_{2N})$  such that  $\phi_f \equiv \phi_s$ , and let  $M_s$  be its linear span. Clearly  $s \in K_s \cap L$ .

Levi and Stark [3] (see also [6]) noted that for a dense Zariski open subset of functions  $s$  in  $L$ , the intersection  $K_s \cap L$  is exactly  $\text{Span}\{s\}$  and may be efficiently reconstructed from the spectral phase  $\phi_s$ . The projection on convex sets (POCS) algorithm developed in [3] applies successive alternate projections on  $K_s$  and  $L$  which eventually converge to  $\text{Span}\{s\}$  from any given initial state  $f_0$ . Let  $\alpha_s$  denote the angle between  $M_s$  and  $L$ . The convergence rate of the POCS algorithm is determined by  $\alpha_s$  as follows: Let  $f_k$  denote the signal obtained after the  $k$ th two-step iteration of the POCS algorithm. It is easy to show that

$$\|f_{k+1} - f_k\| \leq \cos^2(\alpha_s) \|f_k - f_{k-1}\| \tag{1.1}$$

and that  $\cos^2(\alpha_s)$  is the best constant for which (1.1) holds, assuming an arbitrary initial state  $f_0$ .

In a recent paper, Urieli et al. [7] studied the relation between the spatial profile of  $s$  and the angle  $\alpha_s$ . Based on extensive empirical evidence, they conjectured that  $\alpha_s \leq \pi/4$  and that the upper bound is attained when  $s$  has an exponential profile. At the other extreme, they showed that when  $s$  is close to symmetric, the angle  $\alpha_s$  approaches zero (for a symmetric signal the problem is ill-posed).

While our main concern in this article is the behavior of the angle  $\alpha_s$ , it should be noted that in practice the performance of the POCS algorithm is more complex, and is not determined by (1.1) alone. Indeed, a physically realistic model for the amplitude retrieval problem must incorporate model uncertainty and measurement noise. As noted in [1], an exact analysis of noise effects in iterative reconstruction algorithms is difficult. We can however offer the following qualitative remarks concerning the stability of the POCS algorithm with respect to errors: The generic positivity of  $\alpha_s$  combined with the estimate (1.1) guarantees strict contractivity of the POCS, hence convergence to the (essentially unique) solution, for the nominal problem. The continuity of  $\alpha_s$  w.r.t.  $s$  ([7, Theorem 3]) guarantees the same for the perturbed model. To handle measurement noise, one can define at each iteration a ‘‘local angle’’ which dictates the contractivity at that step. The local angle, which is bounded below by  $\alpha_s$ , is continuous w.r.t. the location of the given iterate ([7, Lemma C1]). This guarantees the stability of the convergence process at the presence of relatively small measurement errors. Simulations carried out in [7] appear to support this conclusion.

Returning to the nominal reconstruction problem, our main results (Theorems 1 and 2) are proofs of the above-mentioned conjectures of Urieli et al. [7]. We start with some preliminaries.

The Fourier transform satisfies the Parseval identity  $\langle \widehat{f}, \widehat{g} \rangle = 2N \langle f, g \rangle$ , and the inversion formula  $F^{-1}(g)(m) = (2N)^{-1} \widehat{g}(-m)$ . Note that  $f \in L_{\mathbb{R}}(\mathbb{Z}_{2N})$  iff  $\widehat{f}(-m) = \overline{\widehat{f}(m)}$  for all  $m \in \mathbb{Z}_{2N}$ . The convolution of  $f, g \in L_{\mathbb{C}}(\mathbb{Z}_{2N})$  is given by  $f * g(x) = \sum_{y \in \mathbb{Z}_{2N}} f(y)g(x - y)$  and satisfies  $\widehat{f * g}(x) = \widehat{f}(x)\widehat{g}(x)$ .

Let  $\Lambda_{\mathbb{C}} = \{\lambda \in L_{\mathbb{C}}(\mathbb{Z}_{2N}) : \lambda(j) = \lambda(-j) \text{ for all } j \in \mathbb{Z}_{2N}\}$  and let  $\Lambda_{\mathbb{R}} = \Lambda_{\mathbb{C}} \cap L_{\mathbb{R}}(\mathbb{Z}_{2N})$ . Clearly,  $F(\Lambda_{\mathbb{C}}) = \Lambda_{\mathbb{C}}$ . For a fixed  $s \in L$  let  $W_s = \widehat{s} \cdot \Lambda_{\mathbb{R}} = \{\lambda \widehat{s} : \lambda \in \Lambda_{\mathbb{R}}\}$ , and  $W_s^+ = \{\lambda \widehat{s} : 0 \leq \lambda \in \Lambda_{\mathbb{R}}\}$ . Clearly,  $F^{-1}(W_s^+) = K_s$  and  $F^{-1}(W_s) = M_s$ . Since  $\lambda \widehat{s}(-m) = \overline{\lambda \widehat{s}(m)}$  for all  $\lambda \in \Lambda_{\mathbb{R}}$  and  $m \in \mathbb{Z}_{2N}$ , it follows that  $M_s = F^{-1}(W_s) \subset L_{\mathbb{R}}(\mathbb{Z}_{2N})$  and  $\dim M_s = \dim W_s = |\text{Supp } \widehat{s} \cap \{0, \dots, N\}|$ .

A function  $s \in L$  will be called *generic* if  $M_s \cap L = \text{Span}\{s\}$ . For  $N \leq m \leq 2N - 1$  and  $1 \leq j \leq N$  let

$$D_s(m, j) = \begin{cases} s(m - j) + s(m + j) & \text{if } 1 \leq j \leq N - 1 \\ s(m - N) & \text{if } j = N \end{cases}$$

**Claim 1.**

*If  $s \in L$  satisfies  $\det D_s \neq 0$ , then  $s$  is generic.*

**Proof.** Let  $N \leq m \leq 2N - 1$ . Then for  $\lambda \in \Lambda_{\mathbb{R}}$

$$F^{-1}(\lambda\widehat{s})(m) = \frac{1}{2N}\widehat{\lambda} * s(m) = \frac{1}{2N} \sum_{j=0}^{2N-1} s(m-j)\widehat{\lambda}(j) = \frac{1}{2N} \sum_{j=1}^N D_s(m, j)\widehat{\lambda}(j).$$

Suppose  $D_s$  is non-singular. If  $\lambda \in \Lambda_{\mathbb{R}}$  satisfies  $F^{-1}(\lambda\widehat{s}) \in L$ , then  $F^{-1}(\lambda\widehat{s})(m) = 0$  for all  $N \leq m \leq 2N - 1$ , and hence,  $\widehat{\lambda}(j) = 0$  for all  $1 \leq j \leq N$ . It follows that  $\lambda$  is constant, and therefore  $F^{-1}(\lambda\widehat{s}) \in \text{Span}\{s\}$ .  $\square$

Claim 1 implies that the set of generic functions contains  $S = \{s \in L : \det D_s \neq 0\}$  which is a Zariski open dense subset of  $L$ . Similar genericity criteria appear in [4] and [5]. For  $s \in S$  the angle between  $M_s$  and  $L$  is given by

$$\alpha_s = \min \{\alpha(f, g) : 0 \neq f \in M_s, 0 \neq g \in L, \langle f, s \rangle = \langle g, s \rangle = 0\}$$

where  $\alpha(f, g)$  denotes the angle between  $f$  and  $g$ .

Theorems 1 and 2 were conjectured by Urieli et al. [7] and proved by them for  $N \leq 3$ .

**Theorem 1.**

$\alpha_s \leq \pi/4$  for all  $s \in S$ .

Our main result deals with a case of equality in Theorem 1. A function  $s \in L$  is *geometric with parameter  $q$*  if  $s(j) = q^j$  for all  $0 \leq j \leq N - 1$ .

**Theorem 2.**

If  $s \in L$  is geometric with parameter  $\pm 1 \neq q \in \mathbb{R}$ , then  $s \in S$  and  $\alpha_s = \pi/4$ .

Theorems 1 and 2 are proved in Sections 2 and 3.

## 2. The Upper Bound

For  $t \in \{0, 1\}$  let  $E_t = \{0 \leq j \leq N : j \equiv t \pmod{2}\}$ . We shall need the following technical observation:

**Claim 2.**

Suppose  $0 \neq s \in L$  satisfies  $|\text{Supp } \widehat{s} \cap E_t| \leq 1$  for  $t = 0, 1$ . Then  $N \leq 3$  and  $s$  is a multiple of one of the following functions (written as vectors in  $\mathbb{R}^{2N}$ ):

1.  $N = 2$ :  $(1, 1, 0, 0)$  or  $(1, -1, 0, 0)$ .
2.  $N = 3$ :  $(1, 0, -1, 0, 0, 0)$ .

**Proof.** By assumption there exist  $0 \leq 2k, 2l + 1 \leq N$  and  $\alpha, \beta \in \mathbb{C}$  such that for all  $x \in \mathbb{Z}_{2N}$

$$s(x) = \alpha\omega^{-2kx} + \bar{\alpha}\omega^{2kx} + \beta\omega^{-(2l+1)x} + \bar{\beta}\omega^{(2l+1)x}. \tag{2.1}$$

Let  $0 \leq j \leq N - 1$ , then

$$0 = s(j + N) = \alpha\omega^{-2kj} + \bar{\alpha}\omega^{2kj} - \beta\omega^{-(2l+1)j} - \bar{\beta}\omega^{(2l+1)j}$$

hence, by (2.1)

$$s(j) = 2 \left( \alpha\omega^{-2kj} + \bar{\alpha}\omega^{2kj} \right). \tag{2.2}$$

We consider two cases:

1. If  $2k \equiv 0 \pmod{N}$ , then  $s(j) = 2(\alpha + \bar{\alpha})\epsilon^j$  for a fixed  $\epsilon \in \{\pm 1\}$  and for all  $0 \leq j \leq N - 1$ . It follows that all odd  $m$

$$\widehat{s}(m) = \frac{4(\alpha + \bar{\alpha})}{1 - \epsilon\omega^{-m}} \neq 0.$$

The assumption  $|\text{Supp } \widehat{s} \cap E_1| \leq 1$  then implies that  $N = 2$  and  $s$  is a multiple of  $(1, \epsilon, 0, 0)$ .

2. Suppose  $2k \not\equiv 0 \pmod{N}$ . Let  $2l + 1 \neq m \in E_1$ , then by (2.2)

$$0 = \widehat{s}(m) = \frac{4\alpha}{1 - \omega^{-(2k+m)}} + \frac{4\bar{\alpha}}{1 - \omega^{2k-m}}$$

hence,

$$\left|1 - \omega^{-(2k+m)}\right| = \left|1 - \omega^{2k-m}\right|. \tag{2.3}$$

(2.3) together with  $2k \not\equiv 0 \pmod{N}$  imply  $m = N$ . Since this holds for all odd  $m \neq 2l + 1$  in  $\{0, \dots, N\}$ , it follows that  $N \leq 3$ . The cases  $N = 2$  and  $N = 3, k = 0$  are covered by 1. The remaining possibility is  $N = 3, k = 1$  which implies  $s(j) = 2(\alpha\omega^{-2j} + \bar{\alpha}\omega^{2j})$  for  $0 \leq j \leq 2$ . It can be checked that  $\widehat{s}(1) \neq 0$ ; hence,

$$\widehat{s}(3) = \frac{4\alpha}{1 - \omega} + \frac{4\bar{\alpha}}{1 - \omega^{-1}} = 0.$$

Therefore,  $\alpha/\bar{\alpha} = \omega$  and  $s(j) = 2\bar{\alpha}(\omega^{-2j+1} + \omega^{2j})$  for  $0 \leq j \leq 2$ . It follows that  $s$  is a multiple of  $(1, 0, -1, 0, 0, 0)$ .  $\square$

By checking that  $\det D_s = 0$  for each of the three exceptional cases in Claim 2 we obtain the following

**Corollary 1.**

For any  $s \in S$  there exists a  $t \in \{0, 1\}$  such that  $|\text{Supp } \widehat{s} \cap E_t| \geq 2$ .  $\square$

**Proof of Theorem 1.** Choose  $t \in \{0, 1\}$  such that  $A_t = \text{Supp } \widehat{s} \cap E_t$  satisfies  $|A_t| \geq 2$ . Since

$$\langle \lambda \widehat{s}, \widehat{s} \rangle = \lambda(0) |\widehat{s}(0)|^2 + 2 \sum_{m=1}^{N-1} \lambda(m) |\widehat{s}(m)|^2 + \lambda(N) |\widehat{s}(N)|^2$$

it follows that there exists a  $0 \neq \lambda \in \Lambda_{\mathbb{R}}$  such that  $\text{Supp } \lambda \cap \{0, \dots, N\} \subset A_t$  and  $\langle \lambda \widehat{s}, \widehat{s} \rangle = 0$ . Let  $g = \lambda \widehat{s} \in W_s - \{0\}$ . Since  $g(2j + t + 1) = 0$  for all  $j \in \mathbb{Z}_{2N}$ , it follows that

$$\begin{aligned} \widehat{g}(m + N) &= \sum_{j=0}^{N-1} g(2j + t) \omega^{-(2j+t)(m+N)} \\ &= (-1)^t \sum_{j=0}^{N-1} g(2j + t) \omega^{-(2j+t)m} = (-1)^t \widehat{g}(m). \end{aligned} \tag{2.4}$$

Let  $f = F^{-1}(g) \in M_s - \{0\}$ , and let  $h \in L$  be given by  $h(m) = f(m)$  for  $0 \leq m \leq N - 1$  and  $h(m) = 0$  otherwise. (2.4) implies that  $|f(m)| = |f(m + N)|$  for all  $m \in \mathbb{Z}_{2N}$ , hence,  $\langle h, f \rangle = \langle h, h \rangle = \frac{1}{2} \langle f, f \rangle$ . Let  $\beta$  denote the angle between  $f$  and  $h$ , then

$$\cos \beta = \frac{\langle h, f \rangle}{\|f\| \|h\|} = \frac{1}{\sqrt{2}}.$$

Now

$$\langle h, s \rangle = \langle f, s \rangle = \left\langle F^{-1}(g), F^{-1}(\widehat{s}) \right\rangle = \frac{1}{2N} \langle g, \widehat{s} \rangle = 0$$

hence, both  $f \in M_s - \{0\}$  and  $h \in L - \{0\}$  are orthogonal to  $M_s \cap L$ . It follows that  $\alpha_s \leq \beta = \pi/4$ .

**Remark.** For functions  $s \in S$  which satisfy the additional condition  $\text{Supp } \widehat{s} = \mathbb{Z}_{2N}$  (or equivalently  $\dim M_s = N + 1$ ), Claim 2 and Corollary 1 are of course superfluous.  $\square$

### 3. The Geometric Case

Let  $s \in L$  be geometric with parameter  $q \in \mathbb{R}$ .

**Claim 3.**

$$\det D_s = \begin{cases} (q^N - 1)(q^{2N} - 1)^{\frac{N}{2}-1} & N \text{ even} \\ (q^{2N} - 1)^{\frac{N-1}{2}} & N \text{ odd} \end{cases} \quad (3.1)$$

**Proof.** It is convenient to re-index the rows of  $D_s$  by defining  $C_q(i, j) = D_s(i + N - 1, j)$  for  $1 \leq i, j \leq N$ . More explicitly, let  $A_q, B_q$  be the  $N \times N$  matrices given by

$$\begin{aligned} A_q(i, j) &= \begin{cases} q^{N-1+i-j} & 1 \leq i \leq j \leq N \\ 0 & \text{otherwise} \end{cases} \\ B_q(i, j) &= \begin{cases} q^{i+j-N-1} & N+1 \leq i+j \text{ and } j < N \\ 0 & \text{otherwise} \end{cases} \end{aligned}$$

then  $C_q = A_q + B_q$ .

For  $1 \leq i \leq N$  let  $C_q(i)$  denote the  $i$ th row of  $C_q$ . We prove (3.1) for  $N$  odd: Let  $\zeta = \omega^l$  be any  $2N$ th root of 1. The rows  $\{C_\zeta(i) : 1 \leq i \leq \frac{N+1}{2}\}$  are clearly linearly independent. A routine computation (we omit the details) shows that for  $\frac{N+3}{2} \leq k \leq N$

$$\begin{aligned} C_\zeta(k) &= \zeta^{-(N-1)} \left( C_\zeta(N-k+1) + (\zeta - \zeta^{-1}) \sum_{j=2}^{k-\frac{N+1}{2}} \zeta^{j-2} C_\zeta(N-k+j) \right) \\ &\quad + (\zeta - \zeta^{-1}) (1 + \zeta^{N-1})^{-1} \zeta^{k-\frac{N+3}{2}} C_\zeta\left(\frac{N+1}{2}\right), \end{aligned}$$

therefore,  $C_\zeta(k) \in \text{Span}\{C_\zeta(i) : 1 \leq i \leq \frac{N+1}{2}\}$ . It follows that  $\text{rank } C_\zeta = \frac{N+1}{2}$ , hence, the polynomial  $P(q) = \det C_q$  is divisible by  $(q - \zeta)^{N - \text{rank } C_\zeta} = (q - \omega^l)^{\frac{N-1}{2}}$ . Since  $\deg P(q) = N(N-1)$  it follows that

$$\det D_s = P(q) = \prod_{l=0}^{2N-1} (q - \omega^l)^{\frac{N-1}{2}} = (q^{2N} - 1)^{\frac{N-1}{2}}.$$

The proof of (3.1) for  $N$  even is similar.  $\square$

Assume now that  $q \neq \pm 1$ , then  $s \in S$  by Claim 3. Define an hermitian form  $B$  on  $\Lambda_{\mathbb{C}} \times \Lambda_{\mathbb{C}}$  by

$$B(u, v) = \sum_{m=N}^{2N-1} u * s(m) \overline{v * s(m)} - \sum_{m=0}^{N-1} u * s(m) \overline{v * s(m)}.$$

For  $i \in \mathbb{Z}_{2N}$  let  $e_i \in L_{\mathbb{C}}(\mathbb{Z}_{2N})$  be given by  $e_i(k) = 1$  if  $i = k$  and 0 otherwise. The main ingredient of the proof of Theorem 2 is the following identity.

**Claim 4.**

$$\frac{1 - q^{2N}}{1 - q^2} B(u, v) = B(u, e_N) \overline{B(v, e_N)} - B(u, e_0) \overline{B(v, e_0)}. \quad (3.2)$$

**Proof.** Let  $\theta \in L_{\mathbb{C}}(\mathbb{Z}_{2N})$  be given by

$$\theta(j) = \begin{cases} -1 & \text{if } 0 \leq j \leq N - 1 \\ 1 & \text{if } N \leq j \leq 2N - 1. \end{cases}$$

Note that

$$\widehat{\theta}(m) = \begin{cases} \frac{-4}{1-\omega^{-m}} & m \text{ odd} \\ 0 & m \text{ even} \end{cases}$$

and

$$\widehat{s}(m) = \frac{1 - (-1)^m q^N}{1 - q\omega^{-m}}.$$

Denoting  $A_N = \{(k, l) : 0 \leq k, l \leq 2N - 1 \text{ and } k \not\equiv l \pmod{2}\}$  we obtain

$$\begin{aligned} B(u, v) &= \langle \theta(u * s), v * s \rangle = \frac{1}{2N} \langle (\widehat{\theta}(u * s))^{\widehat{}} , (v * s)^{\widehat{}} \rangle \\ &= \frac{1}{4N^2} \langle \widehat{\theta} * (\widehat{u}s), \widehat{v}s \rangle = \frac{1}{4N^2} \sum_{(k,l) \in A_N} \widehat{\theta}(l - k) \widehat{s}(k) \overline{\widehat{s}(l)} \widehat{u}(k) \overline{\widehat{v}(l)} \\ &= \frac{-1}{4N^2} \sum_{(k,l) \in A_N} \frac{4(1 - q^{2N}) \widehat{u}(k) \overline{\widehat{v}(l)}}{(1 - \omega^{k-l})(1 - q\omega^{-k})(1 - q\omega^l)}. \end{aligned} \tag{3.3}$$

To evaluate the right-hand side of (3.2) we note that  $\widehat{e}_N(m) = (-1)^m$ ,  $B(u, e_0) = -\langle u * s, s \rangle$  and  $B(u, e_N) = \langle u * s, e_N * s \rangle$ . Using the Parseval identity it follows that

$$\begin{aligned} &B(u, e_N) \overline{B(v, e_N)} - B(u, e_0) \overline{B(v, e_0)} \\ &= \langle u * s, e_N * s \rangle \overline{\langle v * s, e_N * s \rangle} - \langle u * s, s \rangle \overline{\langle v * s, s \rangle} \\ &= \frac{1}{4N^2} (\langle \widehat{u}s, \widehat{e}_N s \rangle \overline{\langle \widehat{v}s, \widehat{e}_N s \rangle} - \langle \widehat{u}s, s \rangle \overline{\langle \widehat{v}s, s \rangle}) \\ &= \frac{1}{4N^2} \sum_{k,l=0}^{2N-1} ((-1)^{k+l} - 1) |\widehat{s}(k)|^2 |\widehat{s}(l)|^2 \widehat{u}(k) \overline{\widehat{v}(l)} \\ &= \frac{-1}{4N^2} \sum_{(k,l) \in A_N} \frac{2(1 - q^{2N})^2 \widehat{u}(k) \overline{\widehat{v}(l)}}{|1 - q\omega^{-k}|^2 |1 - q\omega^{-l}|^2}. \end{aligned} \tag{3.4}$$

Now (3.3) and (3.4) imply that (3.2) is equivalent to

$$\begin{aligned} &\frac{2}{1 - q^2} \sum_{(k,l) \in A_N} \frac{\widehat{u}(k) \overline{\widehat{v}(l)}}{(1 - \omega^{k-l})(1 - q\omega^{-k})(1 - q\omega^l)} \\ &= \sum_{(k,l) \in A_N} \frac{\widehat{u}(k) \overline{\widehat{v}(l)}}{|1 - q\omega^{-k}|^2 |1 - q\omega^{-l}|^2}. \end{aligned} \tag{3.5}$$

Since  $\widehat{u}(k) = \widehat{u}(-k)$  and  $\widehat{v}(l) = \widehat{v}(-l)$ , it suffices to show that the sum of the coefficients of  $\widehat{u}(k) \overline{\widehat{v}(l)}$  and  $\widehat{u}(-k) \overline{\widehat{v}(-l)}$  is the same on both sides of (3.5). This follows from the identity:

$$\begin{aligned} &\frac{1}{(1 - \omega^{k-l})(1 - q\omega^{-k})(1 - q\omega^l)} + \frac{1}{(1 - \omega^{l-k})(1 - q\omega^k)(1 - q\omega^{-l})} \\ &= \frac{1 - q^2}{|1 - q\omega^{-k}|^2 |1 - q\omega^{-l}|^2}. \quad \square \end{aligned}$$

Claim 4 now implies the following:

**Claim 5.**

If  $u \in \Lambda_{\mathbb{C}}$  satisfies  $\langle u * s, s \rangle = 0$ , then

$$\sum_{m=N}^{2N-1} |u * s(m)|^2 \geq \sum_{m=0}^{N-1} |u * s(m)|^2.$$

**Proof.** Since  $B(u, e_0) = -\langle u * s, s \rangle = 0$  we obtain by Claim 4:

$$\begin{aligned} & \sum_{m=N}^{2N-1} |u * s(m)|^2 - \sum_{m=0}^{N-1} |u * s(m)|^2 = B(u, u) \\ & = \frac{1 - q^2}{1 - q^{2N}} \left( |B(u, e_N)|^2 - |B(u, e_0)|^2 \right) = \frac{1 - q^2}{1 - q^{2N}} |B(u, e_N)|^2 \geq 0. \quad \square \end{aligned}$$

**Proof of Theorem 2.** By Theorem 1 it suffices to show  $\alpha_s \geq \pi/4$ . Let  $0 \neq f \in M_s$  be orthogonal to  $s$  and let  $0 \neq h \in L$ . Write  $f = \widehat{\lambda} * s$  where  $\lambda \in \Lambda_{\mathbb{R}}$ . Since  $\widehat{\lambda} \in \Lambda_{\mathbb{C}}$ , Claim 5 implies  $\sum_{m=N}^{2N-1} |f(m)|^2 \geq \sum_{m=0}^{N-1} |f(m)|^2$ . Let  $\beta$  denote the angle between  $f$  and  $h$ , then

$$\cos \beta = \frac{|\langle h, f \rangle|}{\|f\| \|h\|} = \frac{\left| \sum_{m=0}^{N-1} f(m)h(m) \right|}{\|f\| \|h\|} \leq \frac{\left( \sum_{m=0}^{N-1} |f(m)|^2 \right)^{\frac{1}{2}} \|h\|}{\|f\| \|h\|} \leq \frac{1}{\sqrt{2}}.$$

It follows that  $\beta \geq \pi/4$ .  $\square$

## References

- [1] Barakat, R. and Newsam, G. (1985). Algorithms for reconstruction of partially known, bandlimited Fourier transform pairs from noisy data, I. the prototypical linear problem, II. the nonlinear problem of phase retrieval, *J. Integral Equations*, **9**, 47–76, 77–125.
- [2] Barakat, R. and Newsam, G. (1985). Algorithms for reconstruction of partially known, bandlimited Fourier transform pairs from noisy data, *J. Opt. Soc. Am.*, **2**, 2027–2039.
- [3] Levi, A. and Stark, H. (1983). Signal restoration from phase by projection into convex sets, *J. Opt. Soc. Am.*, **73**, 810–822.
- [4] Ma, C. (1991). Novel criteria of uniqueness for signal reconstruction from phase, *IEEE Trans. ASSP*, **39**(4), 989–992.
- [5] Oppenheim, A.V. and Lim, J.S. (1981). The importance of phase in signals, *Proc. IEEE*, **69**(5), 529–541.
- [6] Sanz, J.L., Huang, T.S., and Cukierman, F. (1983). Stability of unique Fourier-transform phase reconstruction, *J. Opt. Soc. Am.*, **73**, 1442–1445.
- [7] Urieli, S., Porat, M., and Cohen, N. Optimal reconstruction of images from localized phase, *IEEE Trans. on Image Processing*. In press.

---

Received July 24, 1996  
Revision received March 2, 1998

Mathematics Department, University of Campinas, CP 6065 CEP 13081, Campinas SP, Brazil  
e-mail: nir@ime.unicamp.br

Department of Mathematics, Technion, Haifa 32000, Israel  
e-mail: meshulam@iecor.technion.ac.il